

Bellman-Ford: $O(mn)$

Bellman-Ford(G, s, t)

Initialize $M[i, t] = 0$ for $i=1, \dots, n$ // base case

Initialize $M[0, v] = \infty$ for $v \neq t$ // set distance = ∞ when there is
// no path

for $i=1, \dots, n-1$

 for all $v \in V$

$M[i, v] = M[i-1, t]$ // case 1

 for all $(v, w) \in E$

 if $c_{vw} + M[i-1, w] < M[i, v]$ // case 2 (with first hop w)

$M[i, v] = c_{vw} + M[i-1, w]$

$\text{first}(v) = w$ // to recover the path

 end if

 end for

 end for

end for

return $M[n-1, s]$ // shortest $s-t$ path using at most $n-1$ edges

Note: easy to reduce memory to $O(n)$. See book.